Dibyayan Chakraborty

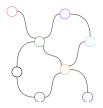
Indian Statistical Institute,Kolkata

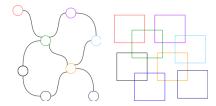
## COCOA 2016

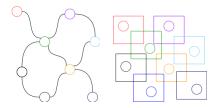
\*\* With Sujoy Bhore, Sandip Das, Sagnik Sen.

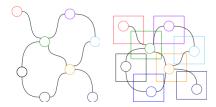


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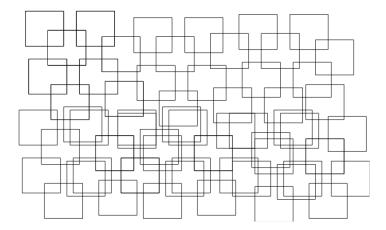


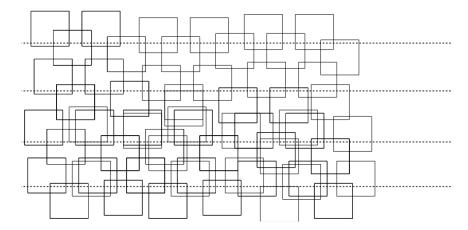
Unfortunately, deciding whether an input graph is a cubicity 2 graph, is NP-hard.

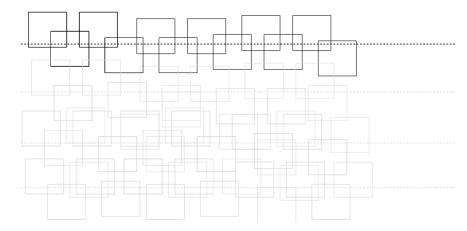
Unfortunately, deciding whether an input graph is a cubicity 2 graph, is NP-hard.

In fact, the complexity of deciding whether an input tree is a cubicity 2 graph, is unknown. (LS Chandran 2014 <sup>1</sup>)

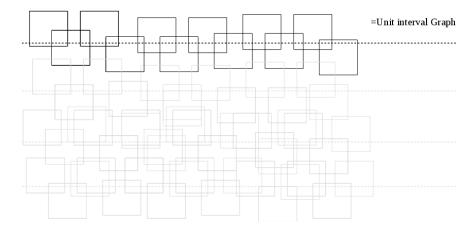
Today we shall solve the above problem "locally".



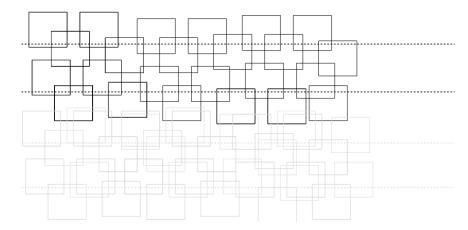




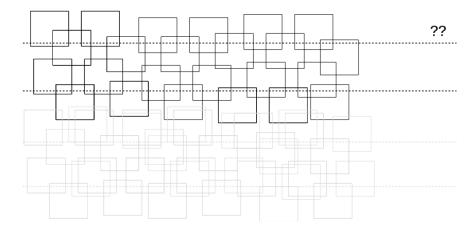
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#### 2SUIG graphs (2-stab unit interval graph )

A graph is a 2SUIG graph if it has intersection representation such that

**Two stab lines in**  $1 + \epsilon$  distance apart.

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- **Two stab lines in**  $1 + \epsilon$  distance apart.
- Axes-parallel unit squares.
- Each unit square intersects exactly one stab line.

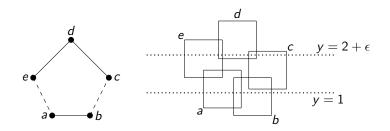


Figure: A representation (right) of a 2SUIG graph (left).

So, let us try to recognize Tree  $\cap$  2SUIG = Tree-2SUIG graphs

In fact we prove that  $\cdots$ 

#### Theorem

There is a O(n) time algorithm to recognize tree-2SUIG graphs.

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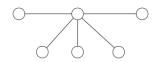
Prove that checking those conditions can be done in linear time.

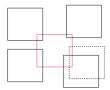
If the input tree satisfies all those conditions, output a 2SUIG representation of the tree.

Otherwise, output a subtree that violets a particular condition.

#### Lemma

If a tree T has a 2SUIG representation then  $\Delta(T) \leq 4$ ,  $\Delta(T)$  is the maximum degree of T.





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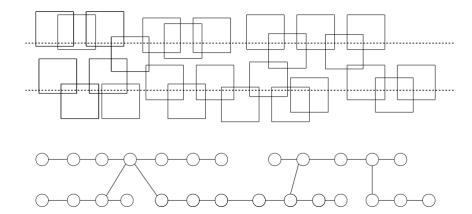


Figure: A typical tree-2SUIG.

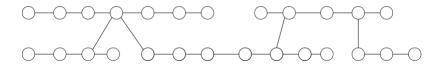
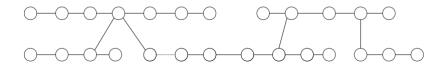
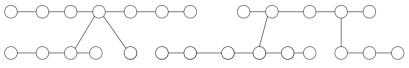


Figure: A typical tree-2SUIG.

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Take an edge and delete it.



Non-unit interval

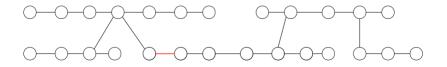
Non-unit interval

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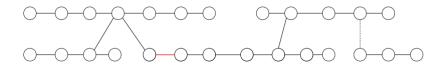
Take an edge and delete it.

If both sub trees are non-unit interval,



Take an edge and delete it.

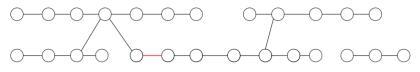
If both sub trees are non-unit interval, call that edge as *red edge*.



Take an edge and delete it.

If both sub trees are non-unit interval, call that edge as *red edge*.

Take another edge and delete it.



Non unit interval

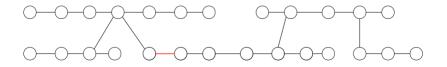
Unit interval

Take an edge and delete it.

If both sub trees are non-unit interval, call that edge as *red edge*.

Take another edge and delete it.

In this case,



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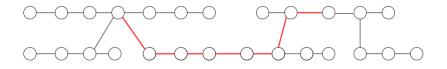
24 / 33

Take an edge and delete it.

If both sub trees are non-unit interval, call that edge as *red edge*.

Take another edge and delete it.

In this case, do nothing.



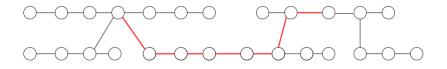
Take an edge and delete it.

If both sub trees are non-unit interval, call that edge as *red edge*.

Take another edge and delete it.

In this case, do nothing.

Get the set of all red edges.



Take an edge and delete it.

If both sub trees are non-unit interval, call that edge as *red edge*.

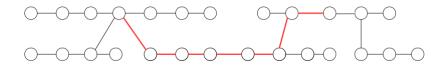
Take another edge and delete it.

In this case, do nothing.

Get the set of all red edges.

#### Lemma

If a tree T has a 2SUIG representation, the set of red edges induces a connected path.

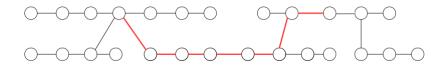


Delete the vertices incident to a red edge.

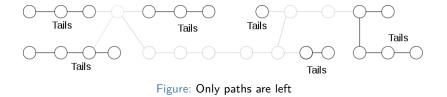


Figure: Only paths are left.

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Delete the vertices incident to a red edge.



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Assuming there are red edges,

Represent the red path appropriately.



Figure: Stretched representation of red path.

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Assuming there are red edges,

Represent the red path (Stretched representation).

Represent the tails appropriately (Shrinked representation).

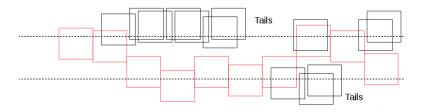


Figure: Shrinked representation of tails.

Length of the tails imposes restriction.

Assuming there are red edges,

Stretched representation of the red path. + Shrinked representation of the tails. U Canonical representation of tree-2SUIG.

Idea is similar, even when there are no red edges.

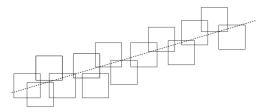
Theorem

There is a O(n) time algorithm to recognize tree-2SUIG graphs.

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#### Open problems

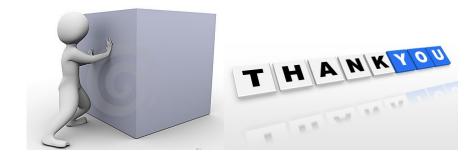
- 1. Determine the complexity of recognizing trees with cubicity 2.
- 2. Recognize those graphs whose intersection representation consists of squares, intersecting a common line, not necessarily axes parallel.



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3. Recognize 2SUIG graphs.

#### Thank You



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